

ON NON-CONSTRUCTIVE THEOREMS OF ANALYSIS AND THE DECISION PROBLEM

R. L. GOODSTEIN

In this note we exhibit a method for strengthening some known results on the impossibility of proving certain classical theorems in recursive analysis. For the concepts and nomenclature of recursive analysis the reader is referred to [3].

We denote by \mathcal{R} some (unspecified) formalisation of recursive arithmetic, and by \mathcal{R}^* an extension of \mathcal{R} to rational numbers and functions, adequate for recursive analysis.

1. If $f(n, x)$ is any rational recursive function, recursively convergent in n , and differentiable in x , relative to n , with relative derivative $f^1(n, x)$ for $0 \leq x \leq 1$, and if

$$f(n, 0) = f(n, 1) = 0$$

then we say that $f(n, x)$ satisfies the conditions of the relative Rolle's theorem and write $f \in RT$.

It was proved in [1] that if $f \in RT$ then there is a recursive ν_k and a recursive c_k such that $n \geq \nu_k \rightarrow f^1(n, c_k) = 0(k)$ is provable in \mathcal{R}^* .

If there exists a recursive sequence c_k , a recursive ν_k and an integer p such that $1/p < c_k < 1 - 1/p$ and

$$(i) \quad n \geq \nu_k \rightarrow f^1(n, c_k) = 0(k)$$

is provable in \mathcal{R}^* (with free variable n) then we say that $f(n, x)$ satisfies the conditions of the uniform Rolle's theorem and write $f \in URT$; if however condition (i) is provable, not necessarily for a variable n , but for each positive integral value of n , then we write $f \in IRT$.

We gave in [2] an example of a function f such that $f \in RT$ but $f \notin URT$. We shall now prove the stronger result that there exists an f such that $f \in RT$ but $f \notin IRT$. We shall in fact show that a proof in \mathcal{R}^* of

$$f \in RT \rightarrow f \in IRT$$

provides a decision method for the class of equations $\varrho(n)=0$, where ϱ is a recursive function which takes only the values 0 and 1, but, as is well known (see [4, pp. 417–418]), this class of equations is undecidable.

2. Given any recursive function $\varrho(n)$ which takes only the values 0 and 1 we define (as in [2])

$$e_0 = 0, \quad e_{n+1} = e_n + \prod_{r=0}^n (1 \div \varrho(r))$$

$$d_0 = 1, \quad d_{n+1} = 1/e_{n+1}$$

and, for $0 \leq x \leq 1$ and $n \geq 3$,

$$f(n, x) = \frac{d_n^4 x (1 \div x)}{d_n^2 + (1 \div 2d_n)x}.$$

It is supposed that $\varrho(0)=0$.

The following properties of these functions are readily provable in \mathcal{R}^* (for details, see [2, pp. 228–230]).

$$(2.1) \quad e_n \leq n.$$

$$(2.2) \quad e_n < n \rightarrow (Er)(r \leq n \ \& \ \varrho(r) = 1).$$

$$(2.3) \quad \text{If } N > n \geq 1 \text{ then } 0 \leq d_n - d_N < 1/n.$$

$$(2.4) \quad \text{For } n \geq 3 \text{ and } 0 \leq x \leq 1, \text{ we have}$$

$$0 \leq f(n, x) \leq d_n^4.$$

$$(2.5) \quad \text{If } 3 \leq n < N \text{ and } 0 \leq x \leq 1 \text{ then}$$

$$0 \leq f(n, x) - f(N, x) < 1/n^4,$$

from which it follows that $f(n, x)$ converges uniformly in x for $0 \leq x \leq 1$.

(2.6) For $0 \leq x \leq 1$, $f(n, x)$ is differentiable in x uniformly in x and n , so that $f(n, x)$ is differentiable in x relative to n , and the relative derivative $f^1(n, x)$ converges uniformly in x , and $f \in RT$.

(2.7) If there is a recursive $V(k)$ such that

$$n \geq V(k) \rightarrow d_n = 0(k)$$

is provable in \mathcal{R}^* for all integers n , then $\varrho(n)=0$ is provable in \mathcal{R} for all integers n .

3. If $f \in IRT$ then (by definition) there exists a recursive $V(k)$, a recursive c_k and an integer p such that $c_k \geq 1/p$ and

$$n \geq V(k) \rightarrow f^1(n, c_k) = 0(k)$$

is provable in \mathcal{R}^* for all integers n .

Since

$$f^1(n, c_k) = \frac{d_n^4(d_n - c_k) \{d_n + (1 - 2d_n)c_k\}}{\{d_n^2 + (1 - 2d_n)c_k\}^2}$$

and

$$\begin{aligned} d_n + (1 - 2d_n)c_k &> d_n^2 + (1 - 2d_n)c_k, \\ d_n^2 + (1 - 2d_n)c_k &\leq (1 - d_n)^2 < 1, \end{aligned}$$

it follows that

$$n \geq V(k) \rightarrow d_n^4(d_n - c_k) = 0(k)$$

is provable in \mathcal{R}^* for all n .

By (2.1), either $d_{p+1} = 1/(p+1)$ or $d_{p+1} > 1/(p+1)$.

If $d_{p+1} = 1/(p+1)$ then $d_n \leq 1/(p+1)$ for $n \geq (p+1)$ so that

$$|d_n - c_k| > 1/[p(p+1)],$$

and therefore

$$n \geq V(4k+p) \rightarrow d_n = 0(k)$$

is provable in \mathcal{R}^* for all n , whence, by (2.7), $\varrho(n) = 0$ is provable in \mathcal{R} for all integers n .

If however $d_{p+1} > 1/(p+1)$ then by (2.2) there is an r between 0 and $p+1$ for which $\varrho(r) = 1$ is provable in \mathcal{R} , and so the hypothesis $f \in IRT$ implies the existence of a decision procedure for the undecidable class of equations $\varrho(n) = 0$.

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UNIVERSITY COLLEGE, LEICESTER, ENGLAND